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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
10/722,021	11/25/2003	Mark Andrew Whittaker Stewart	IS01459MCG 6538	
23330 MOTOROLA,	7590 11/09/2007 INC.	EXAMINER		
LAW DEPARTMENT			LOO, JUVENA W	
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	,		4135	
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			11/09/2007	PAPER

Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

•	Application No.	Applicant(s)				
Office Action Summany	10/722,021	WHITTAKER STEWART, MARK ANDREW				
Office Action Summary	Examiner	Art Unit				
	Juvena W. Loo	4135				
The MAILING DATE of this communication appears on the cover sheet with the correspondence address Period for Reply						
A SHORTENED STATUTORY PERIOD FOR REPLY WHICHEVER IS LONGER, FROM THE MAILING DA - Extensions of time may be available under the provisions of 37 CFR 1.13 after SIX (6) MONTHS from the mailing date of this communication. - If NO period for reply is specified above, the maximum statutory period was reply received by the Office later than three months after the mailing earned patent term adjustment. See 37 CFR 1.704(b).	ATE OF THIS COMMUNICATION 36(a). In no event, however, may a reply be timular apply and will expire SIX (6) MONTHS from a cause the application to become ABANDONE	I. lely filed the mailing date of this communication. D (35 U.S.C. § 133).				
Status						
1) Responsive to communication(s) filed on 11 O	ctober 2007.					
2a)⊠ This action is FINAL . 2b)☐ This	This action is FINAL . 2b) This action is non-final.					
3) Since this application is in condition for allowance except for formal matters, prosecution as to the merits is						
closed in accordance with the practice under E	x parte Quayle, 1935 C.D. 11, 45	33 O.G. 213.				
Disposition of Claims						
4)⊠ Claim(s) <u>1-13 and 15-17</u> is/are pending in the application.						
4a) Of the above claim(s) is/are withdrawn from consideration.						
5) Claim(s) is/are allowed.						
6)⊠ Claim(s) <u>1-13 and 15-17</u> is/are rejected.		·				
7) Claim(s) is/are objected to.						
8) Claim(s) are subject to restriction and/or	r election requirement.					
Application Papers						
9) The specification is objected to by the Examine	r.					
10) The drawing(s) filed on is/are: a) accepted or b) objected to by the Examiner.						
Applicant may not request that any objection to the	drawing(s) be held in abeyance. See	37 CFR 1.85(a).				
Replacement drawing sheet(s) including the correct	ion is required if the drawing(s) is obj	ected to. See 37 CFR 1.121(d).				
11) The oath or declaration is objected to by the Ex	aminer. Note the attached Office	Action or form PTO-152.				
Priority under 35 U.S.C. § 119						
12) Acknowledgment is made of a claim for foreign a) All b) Some * c) None of: 1. Certified copies of the priority documents 2. Certified copies of the priority documents 3. Copies of the certified copies of the prior application from the International Bureau * See the attached detailed Office action for a list of	s have been received. s have been received in Application ity documents have been receive (PCT Rule 17.2(a)).	on No d in this National Stage				
		•				
Attachment(s)						
1) Notice of References Cited (PTO-892)	4) Interview Summary					
2) Notice of Draftsperson's Patent Drawing Review (PTO-948) 3) Information Disclosure Statement(s) (PTO/SB/08) Paper No(s)/Mail Date	Paper No(s)/Mail Da 5) Notice of Informal Pa 6) Other:					

This is in response to amendment filed on October 11, 2007 in which claims 1 to 13,

and 15 - 17 were amended. Claim 14 was cancelled.

Status of Claims

Claims 1–13, and 15-17 are pending, of which claims 1, 8, and 13 are in independent

form.

Claim 8 is rejected under 35 USC 112.

Claims 1 and 3 are rejected under 35 102(b).

Claims 2, 4-13, and 15-17 are rejected under 35 103(a).

1. Applicant's arguments filed October 11, 2007 have been fully considered but they

are not persuasive.

Claim Rejections - 35 USC § 112

2. Claim 8 contains the trademark/trade name InfiniBand. Where a trademark or

trade name is used in a claim as a limitation to identify or describe a particular material

or product, the claim does not comply with the requirements of 35 U.S.C. 112, second

paragraph. See Ex parte Simpson, 218 USPQ 1020 (Bd. App. 1982). The claim scope

is uncertain since the trademark or trade name cannot be used properly to identify any

particular material or product. A trademark or trade name is used to identify a source of

goods, and not the goods themselves. Thus, a trademark or trade name does not

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identify or describe the goods associated with the trademark or trade name. In the present case, the trademark/trade name is used to identify/describe InfiniBand Switch and, accordingly, the identification/description is indefinite.

Double Patenting

3. The nonstatutory double patenting rejection is based on a judicially created doctrine grounded in public policy (a policy reflected in the statute) so as to prevent the unjustified or improper timewise extension of the "right to exclude" granted by a patent and to prevent possible harassment by multiple assignees. A nonstatutory obviousness-type double patenting rejection is appropriate where the conflicting claims are not identical, but at least one examined application claim is not patentably distinct from the reference claim(s) because the examined application claim is either anticipated by, or would have been obvious over, the reference claim(s). See, e.g., *In re Berg*, 140 F.3d 1428, 46 USPQ2d 1226 (Fed. Cir. 1998); *In re Goodman*, 11 F.3d 1046, 29 USPQ2d 2010 (Fed. Cir. 1993); *In re Longi*, 759 F.2d 887, 225 USPQ 645 (Fed. Cir. 1985); *In re Van Ornum*, 686 F.2d 937, 214 USPQ 761 (CCPA 1982); *In re Vogel*, 422 F.2d 438, 164 USPQ 619 (CCPA 1970); and *In re Thorington*, 418 F.2d 528, 163 USPQ 644 (CCPA 1969).

A timely filed terminal disclaimer in compliance with 37 CFR 1.321(c) or 1.321(d) may be used to overcome an actual or provisional rejection based on a nonstatutory double patenting ground provided the conflicting application or patent either is shown to be commonly owned with this application, or claims an invention made as a result of activities undertaken within the scope of a joint research agreement.

Effective January 1, 1994, a registered attorney or agent of record may sign a terminal disclaimer. A terminal disclaimer signed by the assignee must fully comply with 37 CFR 3.73(b).

4. Claims 8, 9, 10, 11, and 12 are provisionally rejected on the ground of nonstatutory obviousness-type double patenting as being unpatentable over claims 5, 6, 7, 8, and 9 of copending Application No. 10/721,213. The conflicting claims are not identical since the current application (No. 10/722,021) is directed to a connection

controller while the other one is directed to an switch (Application No. 10/721,213). The switch and the connection controller have different functional entities and are not patentably distinct from each other because it would have been obvious to one of ordinary skill in the art at the time of the invention to use the controller for the purpose of configuring the switch. The motivation is to enhance the controller's ability to program and update the forwarding instructions in the switch.

This is a <u>provisional</u> obviousness-type double patenting rejection because the conflicting claims have not in fact been patented.

5. Claims 8, 9, 10, 11,12, 13, 15, 16, and 17 are provisionally rejected on the ground of nonstatutory obviousness-type double patenting as being unpatentable over claims 9, 10, 11, 12, 13, 14, 16, 17, and 18 of copending Application No. 10/722,022. The conflicting claims are not identical since the current application (No. 10/722,021) is directed to a connection controller while the other one is directed to a network (Application No. 10/722,022). The connection controller and the network have different functional entities and are not patentably distinct from each other because it would have been obvious to one of ordinary skill in the art at the time of the invention to use the controller for the purpose of calculating an actual traffic pattern through the network. The motivation is to enhance the controller's ability to compute the best path for data transfer.

This is a <u>provisional</u> obviousness-type double patenting rejection because the conflicting claims have not in fact been patented.

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Claim Rejections - 35 USC § 102

6. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that

form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless -

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United

States.

7. Claims 1 and 3 are rejected under 35 USC 102(b) as being anticipated by Bertin

et al. US 6,400,681 B1 (hereinafter Bertin).

Regarding claim 1, Bertin discloses a method, comprising:

a network topology cache coupled to receive network topology data of a network

(Bertin: Figure 5);

a packing algorithm coupled to receive a requested traffic pattern (Bertin: lines

63-65 column 5: when a request for connection is receive, a pre-calculated path is

selected) and to compute an actual traffic pattern of a packet, wherein the packing

algorithm computes an actual traffic pattern using the network topology data and the

requested traffic pattern such that the network operates as a strictly non-interfering

network (Bertin: Figure 7 - all the predetermined paths, corresponds to a specific

destination, are extracted from the Path Table one at a time. If the characteristics of a

path satisfy the request, the path is selected when all the links in the path provide

enough bandwidth. The shortest route, with dedicated bandwidth, is chosen from the selected paths. The predetermined paths are calculated based on the topology information stored in the topology database); and

a logical network state entity coupled to communicate the actual traffic pattern to a source corresponding to the packet (Bertin: lines 53-65 column 11 - connection request is used to reserve bandwidth on every nodes (origin, transit, and destination) in the path. The transit and destination nodes answer the source/request by sending back either a call acceptance or a call reject).

Regarding claim 3, Bertin discloses all the limitations of claim 1 above.

Additionally, Bertin discloses a method in computing an actual traffic pattern comprises:

executing a rearrangement algorithm (Bertin: Figure 7 - all the predetermined paths, corresponds to a specific destination, are extracted from the Path Table one at a time. If the characteristics of a path satisfy the request, the path is selected when all the links in the path provide enough bandwidth. The shortest route, with dedicated bandwidth, is chosen from the selected paths); and

assigning one of a plurality of Destination Location Identifiers (DLID) to the packet such that the network operates as a strictly non-interfering network (Bertin: Figure 4, Path Table (410) – the selected path, with dedicated bandwidth, contains at least the destination node and the selected path).

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Claim Rejections - 35 USC § 103

8. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.

9. Claims 2, 4 - 6, 8, 10 - 13, 16, and 17 are rejected under 35 USC 103(a) as being unpatentable over Bertin et al. (US 6,400,681 B1) in view of Brahmaroutu (US 2003/0033427 A1).

Regarding claim 2, Bertin discloses all the limitations of claim 1 above.

Additionally, Bertin discloses a system that

calculates a plurality of routing trees for a plurality of switches in the network (Bertin: Figure 9, lines 51 column 19 to line 29 column 20; Figure 10, line 32 column 20 to line 4 column 21; and Figure 11, lines 7-33 column 21. The Routing Database and the Topology Database are scanned periodically for identifying new paths or for updating existing ones);

calculates a plurality of Destination Location Identifiers (DLID) and a set of forwarding instructions for each switch (Bertin: line 59 column 12 to line 23 column 13 – Equivalent Capacity of the network connection is first computed. Next, all potential

paths through the network, based on the information stored in the Topology Database, are determined. The algorithm constructs the new potential path by adding one link at a time and ensuring that the bandwidth and quality of services requirements are still met); each of the plurality of DLID corresponds to one of the plurality of routing trees and one of a plurality of destinations (lines 37-39 column 13 – each entry in the table represents a path, between a source node and a destination node, that satisfies specific quality of service and traffic requirements); and

populates a forwarding table (Bertin: line 54 column 18 to line 26 column 19 – paths and their corresponding link information are stored in the Path Table and Link Table respectively. The Path Table and Link Table are both part of Routing Database

Regarding claim 4, Bertin discloses all the limitations of claim 3 above.

Additionally, Bertin discloses a system, wherein

the packet follows a path through at least a portion of a plurality of switches in the network (Bertin: lines 26 - 27, column 7 – incoming data packets are selectively routed onto the outgoing links towards neighboring nodes), and

each of the portion of the switches forwards the packet according to one of the plurality of DLIDs assigned to the packet such that the network operates as a strictly non-interfering network (Bertin: lines 28-29, column 7 – incoming data packet is forwarded according to the routing information contained in the header of the packet; Figure 7 - all the predetermined paths, corresponds to a specific destination, are extracted from the Path Table one at a time. If the characteristics of a path satisfy the

request, the path is selected when all the links in the path provide enough bandwidth.

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The shortest route, with dedicated bandwidth, is chosen from the selected paths. The

predetermined paths are calculated based on the topology information stored in the

topology database).

Regarding claim 5, Bertin discloses all the limitations of claim 4 above.

Additionally, Bertin discloses that each of the portion of the plurality of switches looks up

the one of the plurality of DLIDs assigned to the packet in a forwarding table at each of

the portion of the plurality of switches (lines 28-29 column 7 - routing decisions are

made according to the information contained in the header of the data packets).

Regarding claim 6, Bertin discloses all the limitations of claim 4 above.

Additionally, Bertin discloses that each of the portion of the plurality of switches

forwards the packet in accordance with the one of the plurality of DLIDs assigned to the

packet as in the forwarding table at each the portion of plurality switches (lines 26-27,

column 7 - incoming data packet is forwarded according to the routing information

contained in the header of the packet).

Regarding claim 8, Bertin discloses a method of populating a forwarding table,

comprising:

calculating a plurality of routing trees for a plurality of switches in the network

(Bertin: Figure 9, lines 51 column 19 to line 29 column 20; Figure 10, line 32 column 20

to line 4 column 21; and Figure 11, lines 7-33 column 21. The Routing Database and the Topology Database are scanned periodically for identifying new paths or for updating existing ones);

calculating a plurality of Destination Location Identifiers (DLID) and a set of forwarding instructions for each of the plurality of switches (Bertin: line 59 column 12 to line 23 column 13 – Equivalent Capacity of the network connection is first computed. Next, all potential paths through the network, based on the information stored in the Topology Database, are determined. The algorithm constructs the new potential path by adding one link at a time and ensuring that the bandwidth and quality of services requirements are still met), wherein

each of the plurality of DLIDs corresponds to one of the plurality of routing trees and one of a plurality of end nodes (Bertin: lines 37-39 column 13 – each entry in the table represents a path, between a source node and a destination node, that satisfies specific quality of service and traffic requirements) and wherein

the forwarding instructions create paths appropriate to make the network operate as a strictly non-interfering network (Bertin: Figure 7 - all the predetermined paths, corresponds to a specific destination, are extracted from the Path Table one at a time. If the characteristics of a path satisfy the request, the path is selected when all the links in the path provide enough bandwidth. The shortest route, with dedicated bandwidth, is chosen from the selected paths. The predetermined paths are calculated based on the topology information stored in the topology database); and

populating the forwarding table of each of the plurality of switches in the network with the plurality of DLIDs and the set of forwarding instructions (line 54 column 18 to line 26 column 19 – paths and their corresponding link information are stored in the Path Table and Link Table respectively. The Path Table and Link Table are both part of Routing Database). However, Bertin fails to teach that the switch is an InfiniBand switch.

In the same field of endeavor, Brahmaroutu discloses an InfiniBand switch (Brahmaroutu; Figure 1: an InfiniBand switch fabric) that routes packets based on the information in a forwarding table (Brahmaroutu; Page 3, Section 28 – the packet's header may contain fields such as a destination local identifier used to identify the destination port and data path in the data network, and a source local identifier used to identify the source port used for routing by switches within the network). Thus, it would have been obvious to one of ordinary skill in the art at the time of the invention to apply an InfiniBand switch as disclosed by Brahmaroutu into the network of Bertin. The motivation would have been in reducing latency because InfiniBand switch provides multiple paths between port pairs.

Regarding claim 10, the combinations of Bertin and Brahmaroutu disclose all the limitations of claim 8. Additionally, Bertin discloses each of the plurality of end nodes comprises a destination (Bertin: lines 52-53 column 5 – each node comprises one or more communication devices for receiving or transmitting data packets). However, Bertin fails to teach that the destination is identified by a BaseLID.

In the same field of endeavor, Brahmaroutu discloses every switch and each port may have one or more Local Identifiers (LIDs) (Brahmaroutu; Page 4, Section 31). Thus, it would have been obvious to one of ordinary skill in the art at the time of the invention to apply a local identifier as disclosed by Brahmaroutu into the teaching of Bertin. The motivation would have been in allowing multiple identifiers.

Regarding claim 11, the combinations of Bertin and Brahmaroutu disclose all the limitations of claim 8. Additionally, Bertin discloses that calculating the plurality of routing trees comprises

for each spine node in the network (Bertin: Figure 9, lines 51 column 19 to line 29 column 20; Figure 10, line 32 column 20 to line 4 column 21; and Figure 11, lines 7-33 column 21. The Routing Database and the Topology Database are scanned periodically for identifying new paths or for updating existing ones),

calculating a shortest path from the spine node to each of the plurality of end nodes (Bertin: lines 41 - 44 column 14 - a path is calculated with as few links as possible that supports the quality-of-service requirements of the request).

Regarding claim 12, the combinations of Bertin and Brahmaroutu disclose all the limitations of claim 8. Additionally, Bertin discloses that each of the plurality of routing trees comprises at least a portion of the plurality of switches and the corresponding plurality of links that form a shortest path from one of the plurality of end nodes to a spine node of the network (Bertin: lines 41 - 44, column 14).

packet is created at a source and is addressed to a destination within a network (Bertin:

lines 17 - 26, column 8 - once the optimum paths through the network are calculated,

based on a set of quality of service specifications, so that minimum network resources

are used, the header of the packets is generated in the node), comprising:

executing a rearrangement algorithm for the network (Bertin: Figure 7, all the

predetermined paths, corresponds to a specific destination, are extracted from the Path

Table one at a time. If the characteristics of a path satisfy the request, the path is

selected if its entire links provide enough bandwidth. Moreover, the shortest route, with

dedicated bandwidth, is chosen from the selected paths);

assigning one of a plurality of Destination Location Identifiers (DLID) to the

packet (Bertin: lines 17 - 22 column 8 - the optimum path is put in the header of the

packet); and

the packet following a path through at least a portion of a plurality of switches

from the source to the destination, wherein each of the portion of the plurality of

switches forward the packet according to the one of the plurality of DLIDs assigned to

the packet (Bertin: lines 10-16 column 8 - the packet is routed according to the

information in the header) and wherein

the network operates as a strictly non-interfering network (Bertin: Figure 7 - all

the predetermined paths, corresponds to a specific destination, are extracted from the

Path Table one at a time. If the characteristics of a path satisfy the request, the path is

selected when all the links in the path provide enough bandwidth. The shortest route, with dedicated bandwidth, is chosen from all the selected paths. The predetermined paths are calculated based on the topology information stored in the topology database).

Regarding claim 16, Bertin discloses all the limitations of claim 13 above. Additionally, Bertin discloses that the packet following the path comprises looking up the one of the plurality of DLIDs assigned to the packet in a forwarding table at each of the portion of the plurality of switches along the path form the source to the destination (Bertin: lines 28-29 column 7 – routing decisions are made according to the information contained in the header of the data packets).

Regarding claim 17, Bertin discloses all the limitations of claim 13 above. Additionally, Bertin discloses that the packet following the path comprises each of the portion of the plurality of switches forwarding the packet in accordance with the one of the plurality of DLIDs assigned to the packet as found in a forwarding table at each the portion of the plurality of switches (Bertin: lines 28-29, column 7 – incoming data packet is forwarded according to the routing information contained in the header of the packet).

10. Claim 7 is rejected under 35 USC 103(a) as being unpatentable over Bertin et al. (US 6,400,681 B1) in view of Yang et al. (US Patent 5,940,389).

Regarding claim 7, Bertin discloses all the limitations of claim 1 above. However, Bertin fails to teach that the switch fabric is a CLOS network. In the same field of endeavor, Yang et al. discloses a Benes Network, which is a special case of a CLOS network, can be used as a switch fabric (Yang; lines 32-36 column 6) and that each node in the network has an entry, which is indexed by an identifier and contains information regarding how to transmit received cells to the next node, in the routing table (Yang; lines 30-37 column 2). Thus, it would have been obvious to one of ordinary skill in the art at the time of the invention to apply a CLOS network as disclosed by Yang into the teaching of Bertin. The motivation would have been in reducing latency because CLOS network is well known and expected in the art.

11. Claims 9 and 15 are rejected under 35 USC 103(a) as being unpatentable over Bertin et al., (US 6,400,681 B1) in view of Brahmaroutu (US 2003/0033427 A1) and further in view of Yang et al. (US Patent 5,940,389).

Regarding claim 9, the combination of Bertin and Brahmaroutu disclose all the limitations of claim 8 above. However, Bertin fails to teach that the switch fabric is a CLOS network. In the same field of endeavor, Yang et al. discloses a Benes Network, which is a special case of a CLOS network, can be used as a switch fabric (Yang; lines 32-36 column 6) and that each node in the network has an entry, which is indexed by an identifier and contains information regarding how to transmit received cells to the next

node, in the routing table (Yang; lines 30-37 column 2). Thus, it would have been obvious to one of ordinary skill in the art at the time of the invention to apply a CLOS network as disclosed by Yang into the teaching of Bertin. The motivation would have been in reducing latency because CLOS network is well known and expected in the art.

Regarding claim 15, the combination of Bertin and Brahmaroutu disclose all the limitations of claim 13 above. However, Bertin fails to teach that the switch fabric is a CLOS network. In the same field of endeavor, Yang et al. discloses a Benes Network, which is a special case of a CLOS network, can be used as a switch fabric (Yang; lines 32-36 column 6) and that each node in the network has an entry, which is indexed by an identifier and contains information regarding how to transmit received cells to the next node, in the routing table (Yang; lines 30-37 column 2). Thus, it would have been obvious to one of ordinary skill in the art at the time of the invention to apply a CLOS network as disclosed by Yang into the teaching of Bertin. The motivation would have been in reducing latency because CLOS network is well known and expected in the art.

Response to Argument

Applicant states, as in page 7, that Bertin does not disclose a connection controller that includes a packing algorithm to receive a requested traffic pattern and computes an actual traffic pattern for the packet using the network topology data and

the requested traffic pattern such that the network operates as a strictly non-interfering network. In reply, Bertin discloses that for each connection request, the system selects a pre-calculated path satisfying said connection request in the routing database. However, if no pre-calculated path is found, the system calculates a path that satisfying said connection request (Bertin: line 63 column 5 through line 1 column 6). Bertin further discloses that during the path selection/calculation process, all the information is collected and a Equivalent Capacity of the new network connection is calculated. Potential paths through the network from the origin to the destination node(s) are computed. The path selection/calculation process constructs new potential paths, adds links and nodes on a hop-by-hop basis, and checks at each stage that the user requirements are being met (Bertin: lines 59-67 column 12).

Applicant states, as in page 8, that Bertin does not disclose a process that creates a strictly non-interfering network. In reply, Bertin discloses that for each connection request, the system selects a pre-calculated path satisfying said connection request in the routing database. However, if no pre-calculated path is found, the system calculates a path that satisfying said connection request (Bertin: line 63 column 5 through line 1 column 6). Bertin further discloses that during the path selection/calculation process, all the information is collected and a Equivalent Capacity of the new network connection is calculated. Potential paths through the network from the origin to the destination node(s) are computed. The path selection/calculation process constructs new potential paths, adds links and nodes on a hop-by-hop basis,

and checks at each stage that the user requirements are being met (Bertin: lines 59-67

column 12).

Applicant submits, as in page 9, that Bertin does not disclose each and every

limitation of claim 1 and claim 3. In reply, Bertin discloses a network topology database,

a packing algorithm, and a logical network state entity as explained in the response to

claim 1 above. Bertin also discloses rearrangement algorithm as explained in the

response to claim 3 above.

Applicant submits, as in page 9, that Brahmaroutu also does not discloses a

strictly non-interfering network. In reply, Brahmaroutu discloses the use of multi-path

assignment algorithm to determine and identify all the links that exist between a

destination switch and other switches in the network (Brahmaroutu: Page 7, Section

60). Brahmaroutu also discloses that separate channels may be used for different

functions. In other words, a channel may be used for transporting send request and

reply messages while a separate channel or set of channels may be created for moving

data between two nodes through switches so various sources do not attempt to use the

same network resources at the same time (Brahmaroutu: Page 2, Section 22). In other

words, multiple channels can be dedicated to transport traffic between two nodes.

Applicant submits, as in page 10, the combination of Bertin and Brahmaroutu

does not disclose, teach or otherwise suggest that the forwarding table includes paths

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for a strictly non-interfering network or that the network operates a strictly non-

interfering network as required by independent claims 8 and 13. In reply, Bertin

discloses a path selection/calculation process for creating path through a network that

satisfies user's requirement as explained above. In addition, Brahmaroutu discloses a

technique that allows the use of multi-paths to transport traffic between two network

nodes.

Dependent claims 10-12 and 16-17 remain to be rejected. The Examiner

respectfully submits that claims 10-12 and 16-17 are unpatentable over Bertin in view of

Brahmaroutu for the same reason explained above.

Dependent claims 2 and 4-6 remain to be rejected. The Examiner respectfully

submits that claims 2 and 4-6 are unpatentable over Bertin in view of Brahmaroutu for

the same reason explained above.

Dependent claim 7 remains to be rejected. The Examiner respectfully submits

that claim 7 is unpatentable over Bertin in view of Yang for the same reason explained

above.

Dependent claims 9 and 15 remain to be rejected. The Examiner respectfully

submits that claims 9 and 15 are unpatentable over Bertin in view of Brahmaroutu for

the same reason explained above

Conclusion

13. Applicant's amendment necessitated the new ground(s) of rejection presented in this Office action. Accordingly, **THIS ACTION IS MADE FINAL**. See MPEP § 706.07(a). Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the date of this final action.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Juvena W. Loo whose telephone number is (571) 270-1974. The examiner can normally be reached on Mon.-Thurs: 7:30am - 5:00pm.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Frantz Coby can be reached on (571) 272-4017. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

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Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

Juvena W Loo Examiner Art Unit 4135

SUPERVISORY PATENT EXAMINER